# Priority arguments in descriptive set theory

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## Summary

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- The decomposability conjecture is true, assuming Π<sup>1</sup><sub>2</sub> determinacy. This conjecture characterizes which Borel functions are piecewise continuous.
- ► The proof uses priority arguments in the setting of descriptive set theory. They are carried out using Antonio Montalbán's true stages machinery.

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I was wrong: True stages are a robust framework which have also arisen independently outside of computability theory. True stages are the precise tool needed to deeply understand how membership in a  $\Sigma_n^0$  set is witnessed in boldface descriptive set theory. Montalbán's true stages machinery was essentially invented independently by Louveau and Saint-Raymond in set theory for proving Borel Wadge determinacy in second order arithmetic (Day-Greenberg-Turetsky).

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The Borel hierarchy on a Polish space is defined as follows:

- ▶ The  $\Sigma_1^0$  sets are defined to be the open sets.
- ► The  $\Pi^0_\alpha$  sets are the complements of  $\Sigma^0_\alpha$  sets.
- ▶ A set A is  $\Sigma_{\alpha}^{0}$  if we can express A as a countable union  $A = \bigcup_{i} B_{i}$  where each  $B_{i}$  is a  $\Pi_{\beta}^{0}$  set for some  $\beta < \alpha$

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We compare the complexity of sets by **Wadge reducibility**: if  $A \subseteq X$ , and  $B \subseteq Y$ , then  $A \leq_W B$  if there is a continuous function  $f: X \to Y$  such that for all  $x \in X$ , we have  $x \in A \iff f(x) \in B$ .

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A set B is  $\Sigma_{\alpha}^{0}$ -hard if  $A \leq_{W} B$  for every  $\Sigma_{\alpha}^{0}$  set A. A set is  $\Sigma_{\alpha}^{0}$  complete if it is both  $\Sigma_{\alpha}^{0}$  and  $\Sigma_{\alpha}^{0}$ -hard.

Assume X is Polish.  $A \subseteq X$  is **meager** if its contained in a countable union of nowhere dense sets. So  $B \subseteq X$  is comeager iff there are countably many dense open sets  $\{D_i\}_{i \in \omega}$  so  $\bigcap D_i \subseteq B$ .

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Say a collection  $\{F_n\}_{n\in\omega}$  of closed subsets of X is **good** if  $\{F_n\}_{n\in\omega}$  is closed under finite intersections, and for all  $F_n$  and all open  $U\subseteq X$ , if  $F_n\cap U\neq\emptyset$  then there is some m such that  $F_m\subseteq F_n\cap U$ .

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### Theorem (Abstraction of a folklore technique)

- $A\subseteq X$  is  $\Sigma^0_3$ -hard iff there are good closed sets  $\{F_n\}_{n\in\omega}$  in X so:
  - ▶ If  $x \in X$  is an element of infinitely many  $F_n$ , then  $x \notin A$ .
  - $\triangleright$  For all  $F_n$ , A is comeager in  $F_n$ .

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*Proof sketch:*  $\Leftarrow$ : This is a finite-injury priority argument. Let  $C = \{x \in 2^\omega : (\exists i)(\exists^\infty j) \, x(\langle i,j \rangle) = 1\}$  be the set of reals with some infinite column. This is a complete  $\Sigma_3^0$  set in the Borel hierarchy. It suffices to construct a continuous reduction from C to A. To simplify notation, assume  $X = 2^\omega$ .

#### Proof sketch of $\Leftarrow$ :

For each closed  $F_n$ , let  $\{D_{m,n}\}_{m\in\omega}$  be open dense subsets of  $F_n$  so that  $\bigcap_m D_{m,n}\subseteq A$ .

To each string  $s \in 2^{<\omega}$ , we'll associate a string  $\rho(s) \in 2^{<\omega}$  and a list  $\psi(s) = (F_0^s, \dots, F_{k_s}^s)$  of decreasing closed sets so  $\rho(s) \ge |s|$  and

$$[\rho(s)] \cap \bigcap_{k < k_s} F_k^s \neq \emptyset \tag{*}$$

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For the empty string  $\emptyset$ , let  $\rho(\emptyset) = \emptyset$ , and  $\psi(\emptyset) = \emptyset$ . Define  $\rho(s)$  and  $\psi(s)$  recursively as follows, where  $|s| = \langle i, j \rangle$ :

- For  $s \cap 0$ , extend our list of closed sets). Define  $\psi(s \cap 0)$  extending  $\psi(s)$  so  $|\psi(s \cap 0)| \ge i$  and  $\rho(s \cap 0) \supseteq \rho(s)$  so that (\*) holds for  $s \cap 0$ .
- ▶ (For  $s^1$ , injure closed sets beyond  $F_i^s$  and meet next dense open sets  $D_{m,i}$  in  $F_i^s$ ). Define  $\psi(s^1) = \psi(s) \upharpoonright i$ . Let  $\rho(s^1)$  be such that  $[\rho(s^1)] \subseteq D_{m,i}$  for all  $m \le j$  and (\*) holds.

Define the reduction f from C to A by  $f(x) = \bigcup_n \rho(x \upharpoonright n)$ .

#### Proof sketch of $\Rightarrow$ :

 $\Rightarrow$ : First check that there are such closed sets  $\{F_n\}_{n\in\omega}$  witnessing the theorem for C. If A is  $\Sigma^0_3$  hard, there is a continuous reduction of C to A. By a lemma of Harrington (in Steel (1980) "Analytic sets and Borel isomorphisms"), there is a injective continuous reduction g of C to A. Take image of the sets  $\{F_n\}_{n\in\omega}$  under g.

# Generalizing this theorem throughout the Borel hierarchy

If A is a countable collection of subsets of X, let  $\tau(A)$  denote the topology generated by the subbasis A.

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If A is a countable collection of subsets of X, let  $\tau(A)$  denote the topology generated by the subbasis A.

Given Polish X, say that  $\mathcal{A}_0, \mathcal{A}_1, \ldots, \mathcal{A}_n$  is a **suitable sequence of length** n+1 iff  $\mathcal{A}_0$  is a countable basis of open sets for X,  $\mathcal{A}_m$  is a countable set of  $\Pi_m^0$  subsets of X for  $m \geq 1$ , every  $\mathcal{A}_m$  is closed under finite intersections, and for all m < n,

- 1. If  $B \in \mathcal{A}_0$ , then  $\overline{B} \in \mathcal{A}_1$ , and  $\mathcal{A}_m \subseteq \mathcal{A}_{m+1}$  for m > 0.
- 2. If  $B \in \mathcal{A}_m$ , then  $X \setminus B \in \mathcal{A}_{m+1}$ .
- 3. If  $B \in \mathcal{A}_{m+1}$ , then B is closed in  $\tau(\mathcal{A}_m)$ .
- 4. If  $B \in \mathcal{A}_{m+1}$  and m > 0, then  $\overline{B}^{\mathcal{A}_{m-1}} \in \mathcal{A}_m$ .

Properties (1)-(3) are simple properties which ensure that the topology  $\tau(\mathcal{A}_m)$  is Polish. Property (4) here is the difficult property to satisfy. It is key to the following theorem:

# Characterizing $\Sigma_{n+2}^0$ hardness

#### Theorem (Day-M.)

Suppose X is Polish,  $Y \subseteq X$ , and  $n \ge 1$ . Then Y is  $\Sigma_{n+2}^0$ -hard (i.e. there exists a continuous reduction of a complete  $\Sigma_{n+2}^0$  set to Y) if and only if there exists a closed set  $F \subseteq X$  and a suitable sequence of sets  $A_0, \ldots, A_n$  on F such that

- 1. Y is  $\tau(A_n)$ -meager
- 2. Y is  $\tau(A_{n-1})$ -comeager in A for all  $A \in A_n$

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The proof of  $\Leftarrow$  heavily uses the true stages machinery. We construct a continuous reduction in stages. To each finite string s, we associate an approximation to f(x) which consists of a sequence  $(A_0, A_1, \ldots, A_n)$  of sets with nonempty intersection where  $A_i \in \mathcal{A}_i$ . True stages control the flow of the construction; we will have  $f(x) \in A_i$  if |s| is an i-true stage:

#### Lemma (Montalbán 2014, relativized version)

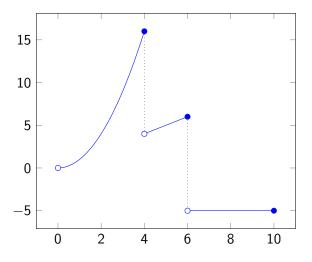
3. If  $\sigma <_{k+1} \tau$ , then  $\sigma <_k \tau$ .

There are partial orders  $\{\leq_k\}_{k\in\omega}$  on  $2^{<\omega}$  and a set  $S_k\subseteq 2^{<\omega}$  such that:

- 1.  $\leq_0$  is the usual prefix ordering:  $\sigma \leq_0 \tau$  iff  $\sigma \subseteq \tau$ .
- 2. The empty sequence  $\emptyset$  has  $\emptyset <_k \sigma$  for every  $\sigma \in T$ .
- 4. If  $\sigma \leq_k \tau$  and  $\sigma \in S_k$ , then  $\tau \in S_k$ .

(\*) If  $\sigma \subseteq \tau \subseteq \rho$  and  $\sigma \leq_{k+1} \rho$  and  $\tau \leq_k \rho$ , then  $\sigma \leq_{k+1} \tau$ . Let  $T_k$  be the tree  $T_k = \{(\sigma_0, \ldots, \sigma_m) : \sigma_i <_k \sigma_{i+1} \land \forall i < k \in I\}$  $m(\neg \exists \tau (\sigma_i <_k \tau <_k \sigma_{i+1}))$  of increasing  $<_k$  sequences. Then for each  $x \in 2^{\omega}$  the restriction of  $T_{k}$  to x:  $T_k \upharpoonright \{(\sigma_0, \ldots, \sigma_m) : (\forall i)\sigma_i \subseteq x\}$  has a single infinite branch, and we say  $\tau \subseteq x$  is a k-true stage of x if  $\tau$  is an element in this unique infinite branch. Then  $\{x : S_k \text{ meets the } k\text{-true stages of } x\}$  is  $\sum_{k=1}^{0}$ complete.

# An application: what functions are piecewise continuous?



Suppose  $f: X \to Y$  is a piecewise continuous function where  $f = \bigcup_{i \in \omega} f_i$ , where the  $f_i$  are partial continuous functions with  $\Delta_n^0$  domains.

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#### Conjecture (2000's, various authors)

 $f: X \to Y$  is a countable union of continuous functions with  $\Delta_n^0$  domains iff the preimage of every  $\Sigma_n^0$  set is  $\Sigma_n^0$ .

#### Conjecture (2000's, the decomposability conjecture)

 $f: X \to Y$  is a countable union of Baire class m functions with  $\Delta_n^0$  domains iff the preimage of every  $\Sigma_{n-m+1}^0$  set is  $\Sigma_n^0$ .

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The decomposability conjecture is true assuming  $\Sigma_2^1$  determinacy.

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These theorems are proved the following way: Suppose  $f: X \to Y$  is not a union of Baire class m functions with  $\Delta_n^0$  domains. Then construct a  $\Sigma_{n-m+1}^0$  set A whose preimage is not  $\Sigma_n^0$  (i.e.  $f^{-1}(A)$  is  $\Pi_n^0$  hard).

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#### Proposition

To prove the decomposability conjecture, it's enough to prove the case where m = n - 1.

Suppose we change our topology  $(X, \tau)$  on X to a new Polish topology  $(X, \eta)$  where we make countably many  $\Pi_n^0$  sets in  $(X, \tau)$  the new basic open sets of  $(X, \eta)$ .

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- ▶ If A is not  $\Sigma_{n+k}^0$  in  $(X,\tau)$ , it is not  $\Sigma_k^0$  in  $(X,\eta)$ .

The converses of these statements are **very** false.

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- ► Then  $f: (X, \eta) \to Y$  is not a union of Baire class 1 functions with  $\Delta_3^0$  domains.
- Apply the techniques of Ding-Kihara-Semmes-Zhao and obtain a  $\Sigma_2^0$  set  $A \subseteq Y$  so that  $f^{-1}(A)$  is not  $\Sigma_3^0$  in  $(X, \eta)$ .

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- Apply the techniques of Ding-Kihara-Semmes-Zhao and obtain a  $\Sigma_2^0$  set  $A \subseteq Y$  so that  $f^{-1}(A)$  is not  $\Sigma_3^0$  in  $(X, \eta)$ .
- ▶ Hope that this set which is not  $\Sigma_3^0$  in  $(X, \eta)$  is not  $\Sigma_n^0$  in the original topology  $(X, \tau)$ .

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If we add this set to our change of topology and try again, we'll eventually succeed at some countable ordinal stage. If not, we would contradict the following:

### Theorem (Harrington 1978, AD)

Fix  $\alpha < \omega_1$ . There is no  $\omega_1$  length sequence of distinct  $\Pi^0_{\alpha}$  sets.

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