The structure of Weihrauch degrees - what we know and what we don't know

Arno Pauly

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MidWest Computability Seminar 2021

2017: The survey

Vasco Brattka, Guido Gherardi & Arno Pauly: Weihrauch Complexity in Computable Analysis. arXiv 1707.03202

And an update

What happened since? What are some interesting open questions?



Arno Pauly:

An update on Weihrauch complexity, and some open questions.

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- Weihrauch reducibility compares multivalued functions between represented spaces.
- ► The induced degrees have a rich algebraic structure.
- Many mathematical theorems can be interpreted as multivalued functions, with the associated Weihrauch degrees measuring the computational content of the theorem.
- The algebraic operations have logic-like meanings regarding such theorems.
- Many concrete theorems have been classified via Weihrauch reducibility; and this classification is reminiscent of reverse mathematics and Brouwerian counterexamples.
- Various techniques have been developed to prove separation results.



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Represented spaces and computability

Definition

A represented space **X** is a pair (X, δ_X) where X is a set and $\delta_X :\subseteq \mathbf{2}^{\mathbb{N}} \to X$ a surjective partial function.

Definition

 $F:\subseteq \mathbf{2}^{\mathbb{N}} \to \mathbf{2}^{\mathbb{N}}$ is a realizer of $f:\subseteq \mathbf{X} \rightrightarrows \mathbf{Y}$, iff $\delta_Y(F(p)) \in f(\delta_X(p))$ for all $p \in \text{dom}(f\delta_X)$.

$$\begin{array}{ccc}
\mathbf{2}^{\mathbb{N}} & \stackrel{F}{\longrightarrow} & \mathbf{2}^{\mathbb{N}} \\
\downarrow \delta_{X} & & \downarrow \delta_{Y} \\
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Definition

 $f:\subseteq X \Rightarrow Y$ is called computable (continuous), iff it has a computable (continuous) realizer.



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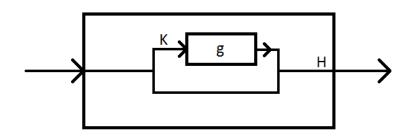
Weihrauch-reducibility

Definition

For $f \subseteq X \Rightarrow Y$, $g \subseteq V \Rightarrow W$ say

$$f \leq_W g$$

iff there are computable $H, K :\subseteq \mathbb{N}^{\mathbb{N}} \to \mathbb{N}^{\mathbb{N}}$, such that $H\langle \operatorname{id}_{\mathbb{N}^{\mathbb{N}}}, GK \rangle$ is a realizer of f for every realizer G of g. \mathfrak{W} denotes the Weihrauch degrees.



Weihrauch reducibility on Baire space

Proposition

For $f,g:\subseteq \mathbb{N}^\mathbb{N} \rightrightarrows \mathbb{N}^\mathbb{N}$ we that $f\leq_W g$ iff there are computable $H,K\subseteq \mathbb{N}^\mathbb{N} \to \mathbb{N}^\mathbb{N}$ with $K:\mathsf{dom}(f)\to \mathsf{dom}(g)$ such that $H(\langle p,q\rangle)\in f(p)$ for all $q\in g(K(p))$.

- Most work on Weihrauch degrees is about classifying specific theorems.
- Then there is work on creating a "scaffolding" of stuff like closed choice principles.
- But only a few papers on the structure of the Weihrauch degrees.
- See http://cca-net.de/publications/weibib.php

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Structures embeddable in the Weihrauch degrees

More algebraic operations

Special subclasses

Some side comments

The big open questions

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Distributive lattice

Theorem (Brattka & Gherardi; Pauly)

The Weihrauch degrees form a distributive lattice;

- with join \sqcup induced by $(f \sqcup g) :\subseteq \mathbf{X} + \mathbf{U} \Rightarrow \mathbf{Y} + \mathbf{U}$, $(f \sqcup g)(0, x) = (0, f(x))$ and $(f \sqcup g)(1, y) = (1, g(y))$,
- ▶ and with meet \sqcap induced by $(f \sqcap g) :\subseteq \mathbf{X} \times \mathbf{U} \Rightarrow \mathbf{Y} + \mathbf{V}$, $(f \sqcap g)(x,y) = (0 \times f(x)) \cup (1 \times g(y))$.

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Special degrees

- ➤ The least element is 0, the trivially true principle without instances.
- ▶ With 1 we denote the degree of $id_{\mathbb{N}^{\mathbb{N}}}$ comprised of all computable problems with a computable instance.
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Theorem (Higuchi & Pauly)

No non-trivial suprema exist in the Weihrauch lattice, meaning either $\sqcup_{i \in \mathbb{N}} f_i$ does not exist, or there is some $N \in \mathbb{N}$ with $\sqcup_{i \in \mathbb{N}} f_i = \sqcup_{i \leq N} f_i$.

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Some non-trivial infima exist, others do not.

Corollary

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Heyting algebra?

Question (Brattka & Gherardi)

Is the Weihrauch lattice a Brouwer algebra, i.e. does

$$\inf_{\leq_W} \{h \mid g \leq_W f \sqcup h\}$$

exist for all f, g?

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The Weihrauch lattice is neither a Brouwer not a Heyting algebra.

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Structures embeddable in the Weihrauch degrees

More algebraic operations

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The big open questions

Medvedev degrees

Definition (Medvedev reducibility)

For $A, B \subseteq \mathbb{N}^{\mathbb{N}}$, $A \leq_M B$ iff $\exists F : B \to A$, F computable. Let \mathfrak{M} denote the Medvedev degrees.

Theorem (Brattka & Gherardi)

 $A\mapsto c_A$, where $c_A(p)=A$, is a meet-semilattice embedding of $\mathfrak M$ into $\mathfrak W$.

Theorem (Higuchi & Pauly)

 $A \mapsto d_A$, where $d_A : A \to \{0\}$, is a lattice embedding of \mathfrak{M}^{op} into \mathfrak{W} . In fact, it is an isomorphism between \mathfrak{M}^{op} and $\{f \in \mathfrak{W} \mid 0 <_W f \leq_W 1\}$.

Question

Is there a lattice-embedding of m into m?



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Is there a lattice-embedding of \mathfrak{M} into \mathfrak{W} ?



Many-one degrees

Definition (Many-one reductions)

For $A, B \subseteq \mathbb{N}$, let $A \leq_m B$ iff there is a computable $F : \mathbb{N} \to \mathbb{N}$ with $F^{-1}(B) = A$.

Theorem (Brattka & Pauly)

The many-one degrees embed into \mathfrak{W} .

Proof.

Let $p,q\in\mathbb{N}^\mathbb{N}$ be Turing incompatible. Map $A\subseteq\mathbb{N}$ to $\chi_A^{p,q}:\mathbb{N} o\{p,q\}$ where $(\chi_A^{p,q})^{-1}(p)=A$.



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We call f join-irreducible, if $f \leq_W g \sqcup h$ implies that $f \leq_W g$ or $f \leq_W h$.

Most "natural" Weihrauch degrees are join-irreducible.

Definition

Let $f \times g : \mathbf{X} \times \mathbf{U} \rightrightarrows \mathbf{Y} \times \mathbf{V}$ be defined via $(y, v) \in (f \times g)(x, u)$ iff $y \in f(x)$ and $v \in g(v)$.

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Sequential composition

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Let
$$f\star g=\sup_{\leq_{\mathrm{W}}}\{F\circ G\mid F\leq_{\mathrm{W}}f\wedge G\leq_{\mathrm{W}}g\}.$$

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Theorem (Dzhafarov, Goh, Hirschfeldt, Patey & Pauly) $RT_2^2 \leq_W SRT_2^2 \star COH$, but RT_2^2 and $SRT_2^2 \times COH$ are incomparable.

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Closure under composition

Definition (Neumann & Pauly)

An input for f^{\diamond} is a description of an abstract register machine operating on represented spaces with computable functions and f as operations, together with an input on which the register machine halts. The output is whatever the register machine outputs.

This is *supposed* to capture closure under composition.

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Proposition

 f^* is the least Weihrauch degree above f satisfying $1 \leq_W f^*$ and $f^* \times f^* \equiv_W f^*$.

Theorem (Westrick 2020)

- ► Open since CCA 2015
- There is a constant function f and a multivalued function g such that $f \leq_W g^{\diamond}$, but no fixed finite number of applications of g suffices

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- 2. $f \sqcup g$, letting us choose between f and g (AND)
- 3. $f \times g$, letting us both f and g in parallel (AND)
- 4. $f \star g$, letting us first use g, then f (AND)
- 5. $f \rightarrow g = \min\{h \mid g \leq_W f \star h\}$ (Implication)
- 6. f^* , f^{\diamond} letting us use f finitely many times, in parallel or consecutively (bang, bang)
- 7. \hat{f} , letting us use f countably many times in parallel (bang)
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- 8. (and more)



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The idea

Sometimes, we can understand a Weihrauch degree by figuring out how it relates to "simple" Weihrauch degrees.

Definition (Dzhafarov, Solomon & Yokoyama)

Let the first-order part of a Weihrauch degree *f* be:

$${}^{1}f := \sup_{\leq_{\mathrm{W}}} \{g : \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N} \mid g \leq_{\mathrm{W}} f\}$$

Definition (Valenti, Goh & Pauly)

Fix a represented space \mathbf{X} . The deterministic part of a Weihrauch degree f is

$$\mathsf{Det}_{\mathbf{X}}(f) := \sup_{\leq_{\mathsf{W}}} \{g : \subseteq \mathbb{N}^{\mathbb{N}} \to \mathbf{X} \mid g \leq_{\mathsf{W}} f\}$$

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Some questions and results

Proposition (Hoyrup)

There is an f with $\mathsf{Det}_{\mathbb{N}^{\mathbb{N}}}(f) <_W \mathsf{Det}_{\mathbb{R}}(f)$.

Proposition (de Brecht, Pauly & Schröder)

For overt choice $\mathbf{VC}_{\mathbb{Q}} :\subseteq \mathcal{V}(\mathbb{Q}) \rightrightarrows \mathbb{Q}$ it holds that ${}^{1}(\mathbf{VC}_{\mathbb{Q}}) \equiv_{W} \mathrm{Det}_{\mathbb{N}^{\mathbb{N}}}(\mathbf{VC}_{\mathbb{Q}}) \equiv_{W} 1$, but $\mathbf{VC}_{\mathbb{Q}}$ is not computable.

Question (Valenti, Goh & Pauly)

Is there some f with $\mathsf{Det}_{\mathbb{N}}(f) <_W \mathsf{Det}_{\mathbb{N}^{\mathbb{N}}}(^1f)$? (It always holds that $\mathsf{Det}_{\mathbb{N}}(f) \equiv_W ^1 \mathsf{Det}_{\mathbb{N}^{\mathbb{N}}}(f)$)

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Irreducibility

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There are $f, g <_W \lim with f \times g \equiv_W \lim$.

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- Every countable distributive lattice embeds into the Weihrauch degrees (via the Medvedev degrees).
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If we relativize Weihrauch reducibility relative to an arbitrary oracle, we get continuous Weihrauch reducibility.

Question

How do the Weihrauch degrees inside a given continuous Weihrauch degree look like?

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